



Middlebury

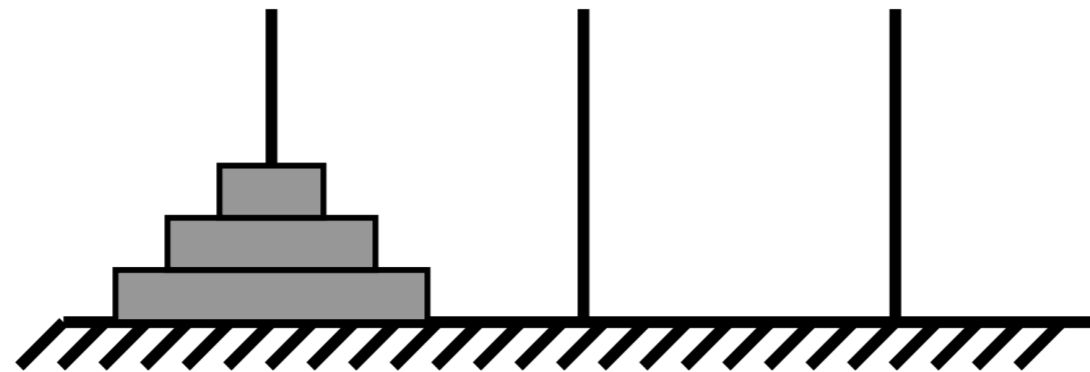
CSCI 201: Data Structures

Spring 2025

Lecture 5M: Recursion

Goals for today:

- Associate a call to a method with a **frame**.
- Use **recursion** to solve problems by building a solution from solutions to smaller problems.
- Ensure we always have a **base case** in our recursive solutions.
- Use **tail recursion** so that there are no pending operations after the recursive call.
- Practice some more with big-O notation.

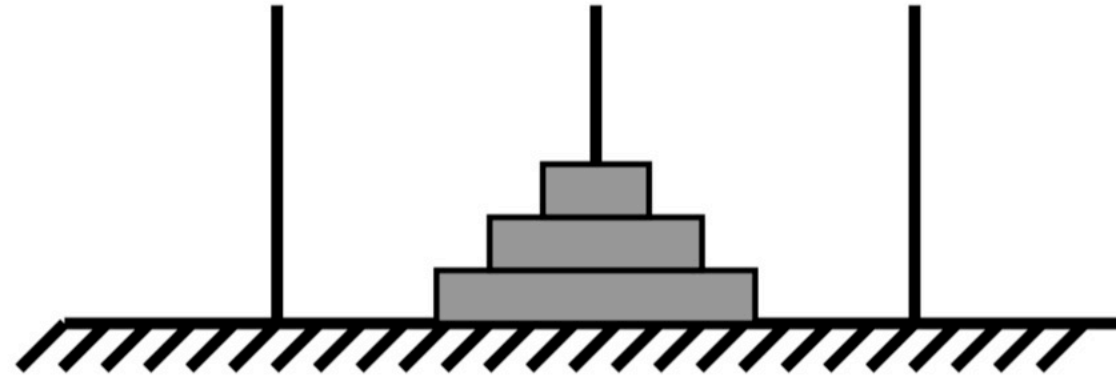


Rules: (also play here! <https://www.mathsisfun.com/games/towerofhanoi.html>)

1. Move one disk at a time.
2. Every move displaces a disk from the top of one stack to the top of another (or empty rod).
3. No larger disk can ever be placed on top of a smaller one.

Counting the number of "moves" in the Tower of Hanoi problem.

"Moves"
operations
 $T(n)$



n	# moves
1	1
2	3
3	7
4	15
5	31

$$T(1) = 1$$

$$\begin{aligned} T(n) &= 2T(n-1) + 1 \\ &= 2(2T(n-2) + 1) + 1 \\ &\dots \\ &= 2^n - 1 \quad O(2^n) \end{aligned}$$

↑
 $2^n - 1$

Two ingredients for a recursive method:

- **Base case:** represents a problem you know how to solve.
- **Recursive case:** builds the solution to a problem from smaller subproblems.

Arithmetic series: $1 + 2 + 3 + \dots + (n - 1) + n$

$= \frac{n(n+1)}{2}$

Factorial: $1 \cdot 2 \cdot 3 \cdot \dots \cdot (n - 1) \cdot n$

```
1 public static int sum(int n) {
2     if (n == 1) { // base case
3         return 1;
4     }
5     // recursive case
6     return n + sum(n - 1);
7 }
```

approach $n == 1$

```
1 public static int factorial(int n) {
2     if (n < 2) { // base case
3         return 1;
4     }
5     // recursive case
6     return n * factorial(n - 1);
7 }
```

approach $n < 2$

The recursive case needs to make progress towards the base case!

What does the **Call Stack** look like?

Open **GoBackExample.java** and add breakpoint on Line 6.

RUN AND DEBUG

No Configurations

VARIABLES

Local

n = 1

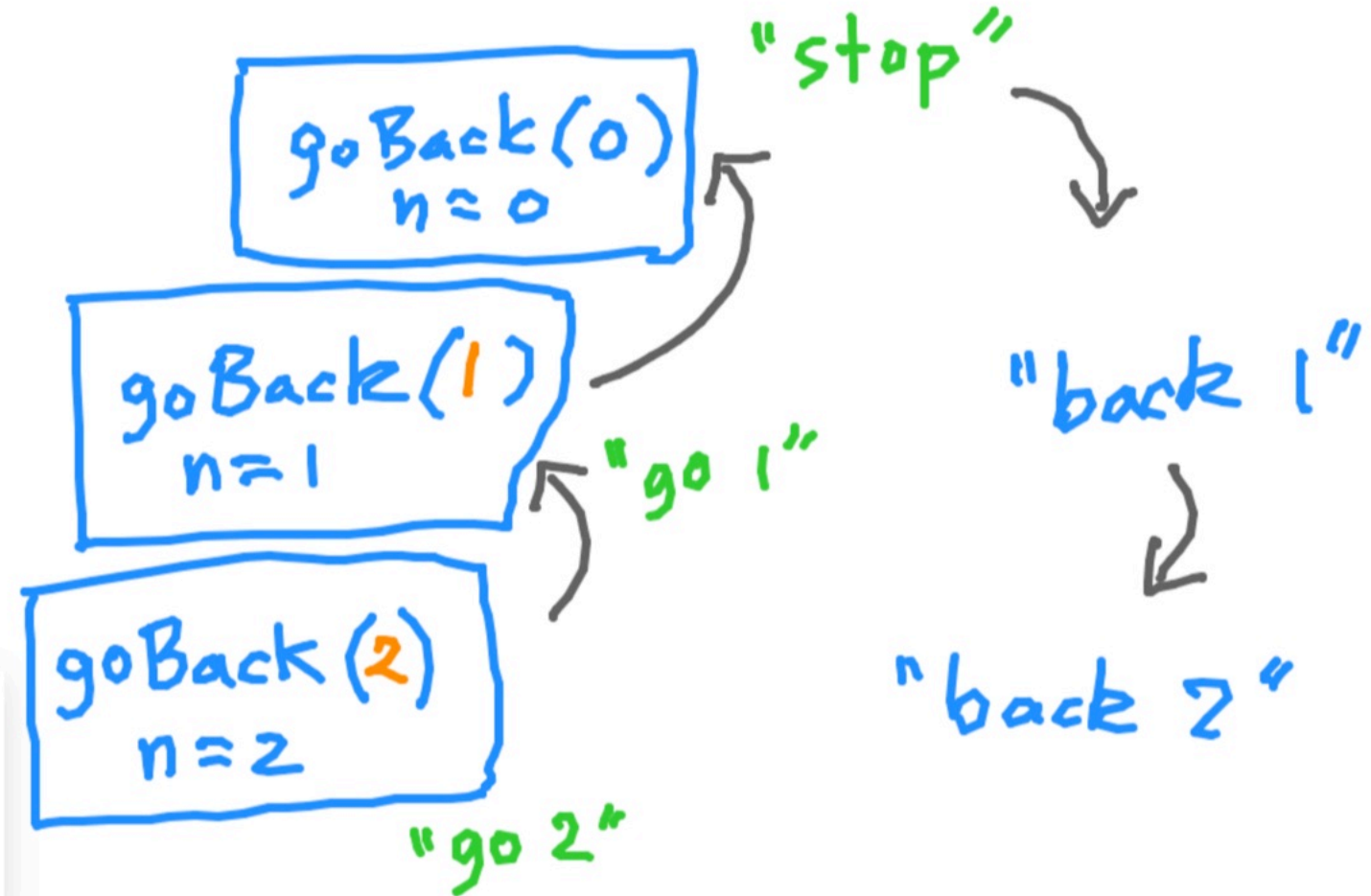
WATCH

CALL STACK

Thread [main] PAUSED ON STEP

GoBackExample.goBack(int)	GoBackExample.java	10:1
GoBackExample.goBack(int)	GoBackExample.java	9:1
GoBackExample.goBack(int)	GoBackExample.java	9:1
GoBackExample.goBack(int)	GoBackExample.java	9:1
GoBackExample.goBack(int)	GoBackExample.java	9:1
GoBackExample.goBack(int)	GoBackExample.java	9:1
GoBackExample.main(String[])	GoBackExample.java	14:1

```
1 public static void goBack(int n) {
2     if (n == 0) {
3         System.out.println("Stop");
4     } else {
5         System.out.println("Go " + n);
6         goBack(n - 1);
7         System.out.println("Back " + n);
8     }
9 }
```



recursion (n.)

"return, backward movement," 1610s, from Latin *recursionem* (nominative *recursio*) "a running backward, return," noun of action from past-participle stem of *recurrere* "run back" (see **recur**).



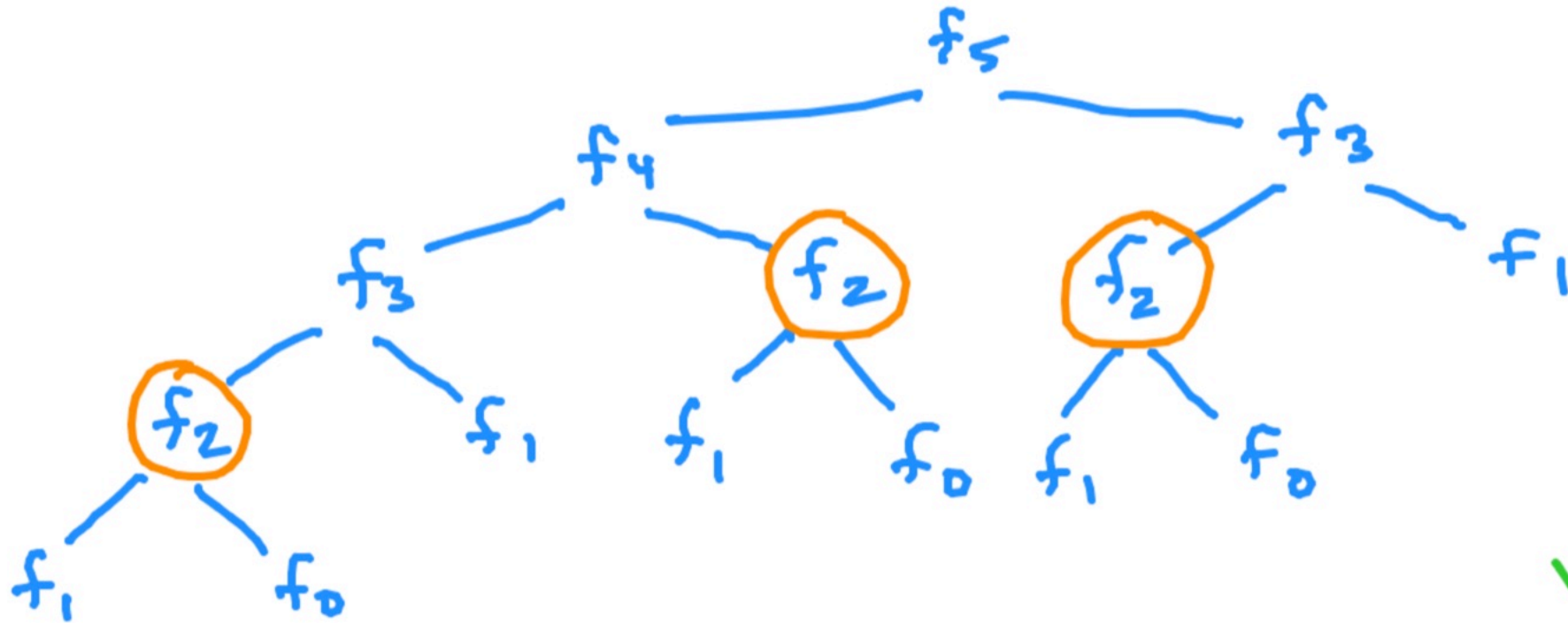
Exercise: write a recursive function to compute Fibonacci numbers.

$$f(n) = f(n - 1) + f(n - 2), \quad \text{with } f(0) = 0, f(1) = 1.$$

```
1 public static int fib(int n) {  
2     if (n < 2) {  
3         return n;  
4     }  
5     return fib(n - 1) + fib(n - 2);  
6 }
```

levels = n
branches = 2

Try $n = 40, n = 45, n = 48$. What do you notice?

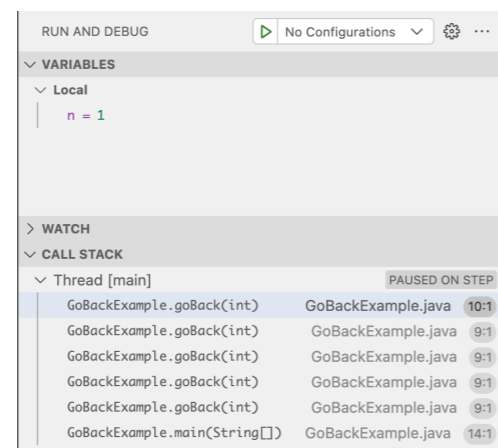


$\approx 2^n$ calls
to fib
 $O(2^n)$
exponential
very slow!

The downsides of recursion.

- Performance might not be great.
- Recursion depth is limited by the stack size.

Go back to **GoBackExample.java** (remove breakpoints) and try $n = 100000$.



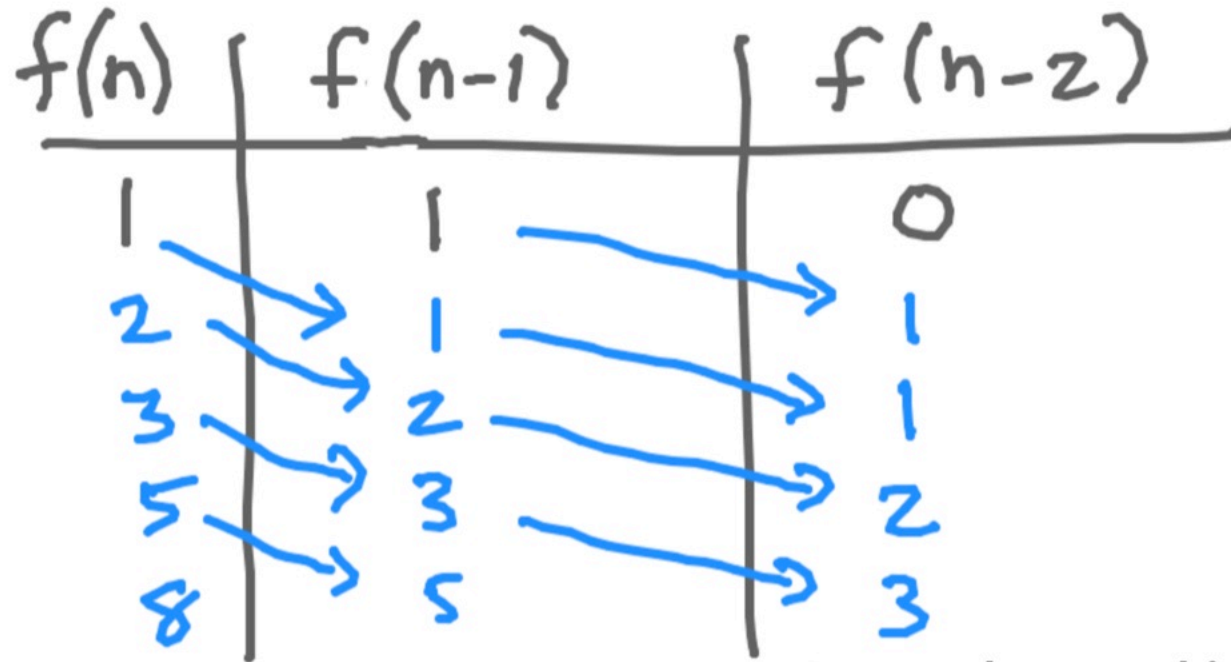
```
1 public static void goBack(int n) {  
2     if (n == 0) {  
3         System.out.println("Stop");  
4     } else {  
5         System.out.println("Go " + n);  
6         goBack(n - 1);  
7         System.out.println("Back " + n);  
8     }  
9 }
```

```
Go  
Go  
Go  
Exception in thread "main" java.lang.StackOverflowError  
at java.base/java.io.FileOutputStream.write(FileOutputStream.java:349)  
at java.base/java.io.BufferedOutputStream.flushBuffer(BufferedOutputStream.java:81)  
at java.base/java.io.BufferedOutputStream.flush(BufferedOutputStream.java:142)  
at java.base/java.io.PrintStream.write(PrintStream.java:570)  
at java.base/sun.nio.cs.StreamEncoder.writeBytes(StreamEncoder.java:234)  
at java.base/sun.nio.cs.StreamEncoder.implFlushBuffer(StreamEncoder.java:313)  
at java.base/sun.nio.cs.StreamEncoder.flushBuffer(StreamEncoder.java:111)  
at java.base/java.io.OutputStreamWriter.flushBuffer(OutputStreamWriter.java:178)  
at java.base/java.io.PrintStream.writeln(PrintStream.java:723)  
at java.base/java.io.PrintStream.println(PrintStream.java:1028)  
at GoBackExample.goBack(GoBackExample.java:8)  
at GoBackExample.goBack(GoBackExample.java:9)  
at GoBackExample.goBack(GoBackExample.java:9)
```



Can we make our Fibonacci number calculation more efficient?

Let's try to write it as a **while**-loop instead.



```
1 public static int fibWhile(int n) {
2     int fnMinus2 = 0;
3     int fnMinus1 = 1;
4     int i = n;
5     while (i > 1) {
6         int fn = fnMinus1 + fnMinus2;
7         fnMinus2 = fnMinus1;
8         fnMinus1 = fn;
9         i--;
10    }
11    return fnMinus1;
12 }
```

Can we do something similar using recursion?

```
1 public static int fib(int n) {
2     if (n < 2) return n;
3     return fibHelper(n, 1, 0);
4 }
5
6 private static int fibHelper(int n, int fnMinus1, int fnMinus2) {
7     if (n == 1) return fnMinus1;
8     return fibHelper(n - 1, fnMinus1, fnMinus2);
9 }
```


Can we make our *recursive* Fibonacci number calculation more efficient?

```
1 public static int fibWhile(int n) {
2     int fnMinus2 = 0;
3     int fnMinus1 = 1;
4     int i = n;
5     while (i > 1) {
6         int fn = fnMinus1 + fnMinus2;
7         fnMinus2 = fnMinus1;
8         fnMinus1 = fn;
9         i--;
10    }
11    return fnMinus1;
12 }
```

```
1 public static int fib(int n) {
2     if (n < 2) return n;
3     return fibHelper(n, 1, 0);
4 }
5
6 private static int fibHelper(int n, int fnMinus1, int fnMinus2) {
7     if (n == 1) return fnMinus1;
8     return fibHelper(n - 1, ↑ , ← );
9 }
```

fnMinus1 + fnMinus2

fnMinus1

Fill in the blanks! See [slido.com # 4017686](https://www.slido.com/#4017686)

☰ CS 201 Lecture 9



☰ What should the recursive call to fibHelper look like on Line 8?

17 👤

fibHelper(n - 1, fnMinus1, fnMinus1 + fnMinus2);

fibHelper(n - 1, fnMinus2, fnMinus1 + fnMinus2);

fibHelper(n - 1, fnMinus1 + fnMinus2, fnMinus1);

Send

Voting as Anonymous



Two things going on here:

- Computing any given Fibonacci number *once*.
- **Tail-Recursive**: the recursive call is the **last** thing done in the method.
 - **Strategy to develop**: think about how to write it as a **while**-loop first.

helper

[

```
1 public static int fib(int n) {  
2     if (n < 2) return n;  
3     return fibHelper(n, 1, 0);  
4 }  
5  
6 private static int fibHelper(int n, int fnMinus1, int fnMinus2) {  
7     if (n == 1) return fnMinus1;  
8     return fibHelper(n - 1, fnMinus1 + fnMinus2, fnMinus1);  
9 }
```

How many times is **fibHelper** called?

same as while loop
 $O(n)$

why tail recursion?

→ some compilers
detect and will

optimize creation of stack frames

Exercise: rewrite **factorial** to be tail-recursive.

```
int result = 1;
int i = n;
while (i > 1) {
    result *= i;
    i--;
}
```

```
1 public static int factorial(int n) {
2     if (n < 2) { // base case
3         return n;
4     }
5     // recursive case
6     return n * factorial(n - 1);
7 }
```

```
1 public static int factorialTail(int n) {
2     return factorialTailHelper(n, 1);
3 }
4
5 private static int factorialTailHelper(int n, int result) {
6     if (n == 1) return result;
7     return factorialTailHelper(n - 1, n * result);
8 }
```

In this case, we still have $\mathcal{O}(n)$ multiplications ($*$), but a compiler might perform *tail call optimization* (TCO), eliminating the need for a new stack frame (just not in **Java**).

Back to **GoBackExample.java**.

Try to pinpoint *exactly* the value of **n** that gives a stack overflow.

Design your own algorithm!

This is the algorithm *binary search*.

```
1 public static int binarySearch(int value, int[] sortedArray) {
2     return binarySearch(value, sortedArray, 0, sortedArray.length - 1);
3 }
4
5 private static int binarySearch(int value, int[] sortedArray, int left, int right) {
6     if (right < left) {
7         return -1;
8     }
9
10    int mid = left + (right - left) / 2;
11    if (sortedArray[mid] == value) {
12        return mid;
13    } else if (sortedArray[mid] < value) {
14        return binarySearch(value, sortedArray, mid + 1, right); // search right half
15    } else {
16        return binarySearch(value, sortedArray, left, mid - 1); // search left half
17    }
18 }
```


More practice

Convert the following (recursive) string reversal method to use tail recursion.

```
1 public static String reverse(String message) {  
2     if (message.length() < 2) return message;  
3     return message.charAt(message.length() - 1) + reverse(message.substring(0, message.length() - 1));  
4 }
```

```
1 public static String reverseTail(String message) {  
2     return reverseHelper(message, "");  
3 }  
4  
5 private static String reverseHelper(String message, String result) {  
6     if (message.length() == 0) {  
7         return result;  
8     }  
9     int end = message.length() - 1;  
10    return reverseHelper(message.substring(0, end), result + message.charAt(end));  
11 }
```

See you on Wednesday!

- [Lab 4](#) (Bucket Sort) due tonight at 11:59pm.
- Work on [Homework 4](#)!
 1. Implement Radix Sort.
 2. Derive number of `=` for `DIYList add` method with a different growth technique.
- Reminder that Noah ([go/noah](#)) and Smith ([go/smith](#)) have office hours throughout the week and the 201 Course Assistants have drop-in hours in the late afternoons/evenings ([go/cshelp](#)).